

Introduction to higher topoi

2: Descent into ∞ -topoi

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Grothendieck topos: category \mathcal{E} such that \exists

(1) small ^{1.}category \mathcal{C}

(2) fully faithful $\mathcal{E} \xrightarrow{i} \mathcal{P}(\mathcal{C}) := \text{Fun}(\mathcal{C}^{\text{op}}, \text{Set})$, with

(3) a left adjoint $l \dashv i$, s.t.

(4) l preserves finite limits (left exact)

Correspondence:

(lex localizations of
 $\text{Fun}(C^{\text{op}}, \text{Set})$)



(Grothendieck
topologies on C)

"well-known"

Basic properties of Gr. Topoi: follow immediately

• Cartesian closure: $\Pi(X, Y) \in \mathcal{P}(C) \xrightarrow{l} \Sigma$
 $\Upsilon \in \Sigma \Rightarrow \Pi(X, \Upsilon) \in \Sigma$ $\begin{matrix} l(X=Y) \\ \text{"} \\ lX=lY \end{matrix}$

• Slices: $\mathcal{P}(C) \xrightleftharpoons[i]{l} \Sigma \Rightarrow \mathcal{P}(C)_{/X} \xrightleftharpoons[i]{l} \Sigma_{/X}$
 $\bigcup_{X \in \mathcal{L}X} \mathcal{P}(C/X)$ - lex

• Subobject classifier: $\Omega \in \mathcal{P}(C) \xrightarrow{l} \Sigma$
 $l \Rightarrow \Omega \xrightarrow{j} \Omega \text{ idempotent} \Rightarrow j\Omega \in \Sigma$ ✓

∞ -topos : ∞ -category \mathcal{E} such that \exists

(1) small ∞ -category \mathcal{C}

(2) **accessible** + fully faithful $\mathcal{E} \xrightarrow{i} \text{PSh}(\mathcal{C}) = \text{Fun}(\mathcal{C}^{\text{op}}, \mathbb{S})$, with
 \uparrow presheaf of ∞ -groupoids

(3) a left adjoint $l \dashv i$, s.t.

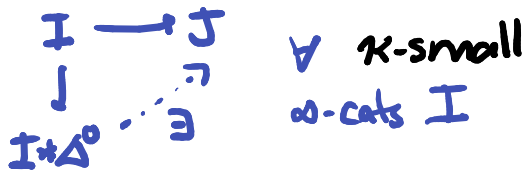
(4) l preserves finite limits ("left exact")

"left exact localization of presheaves"

Accessible functor $F: C \rightarrow D$ of ∞ -categories:

- C, D cocomplete
- F preserves κ -filtered colimits
(some regular cardinal κ)

• J is κ -filtered if



In Gr-topos accessibility
of j follows from
other properties

Lurie, "HTT", Ch. 5 \Leftarrow

Presentable ∞ -category: \mathcal{A} such that \exists

(1) small ∞ -category \mathcal{C}

(2) **accessible** + fully faithful $\mathcal{A} \xrightarrow{i} \text{PSh}(\mathcal{C}) = \text{Fun}(\mathcal{C}^{\text{op}}, \mathcal{S})$, with

(3) a left adjoint $l \dashv i$

∞ -cat generalization of "local presentable categories"

Lurie, "HTT", Ch 5. [also Simpson (arXiv 1999)]

What about Grothendieck topologies?

• Given (C, τ) (Grothendieck site on a 1-category C)

$$\Rightarrow \text{Sh}(C, \tau) \subseteq \text{PSh}(C) := \text{Fun}(C^{\text{op}}, \mathcal{S})$$

∞-topos sheaves

Ex: $X = \text{top'l space} \Rightarrow \text{Sh}(X) \subseteq \text{Fun}(\text{Open}_X^{\text{op}}, \mathcal{S})$

defined in previous hour

Characterize Grothendieck topoi

Giraud Theorem: A 1-category \mathcal{E} is a

Grothendieck topos iff

- \mathcal{E} is locally presentable.
- colimits are universal
- coproducts are disjoint
- equivalence relations are effective.

Characterize ∞ -topoi:

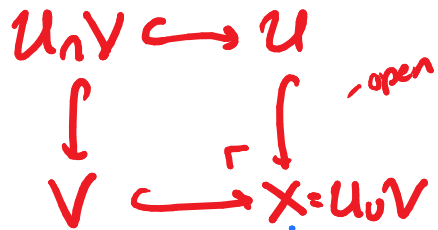
Theorem: An ∞ -category \mathcal{E} is an ∞ -topos iff

- \mathcal{E} is presentable
- colimits are universal.
- colimits satisfy descent \Rightarrow

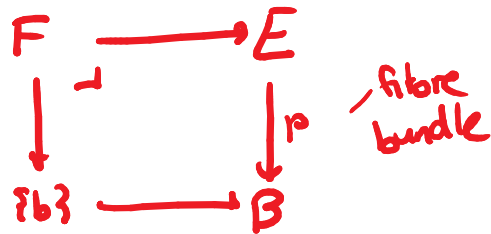
"descent" }

} have roots in homotopy

Homotopy limits/colimits : in top'l spaces



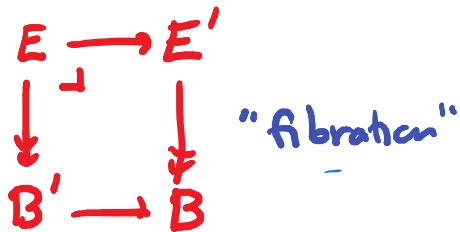
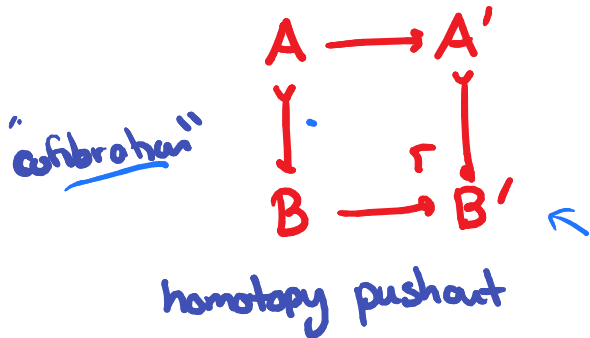
\Rightarrow behaves well wrt H_*



\Rightarrow Behaves well wrt π_*

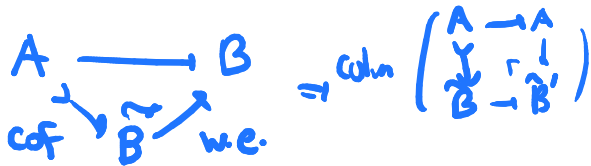
Homotopy pushouts/pullbacks:

in Top or sSet
[or (proper) QMC]

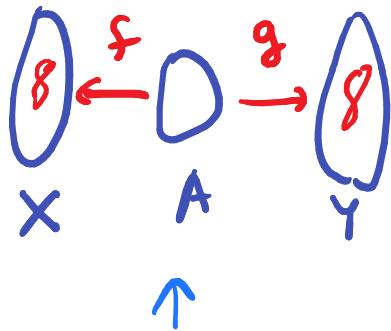


homotopy pullback

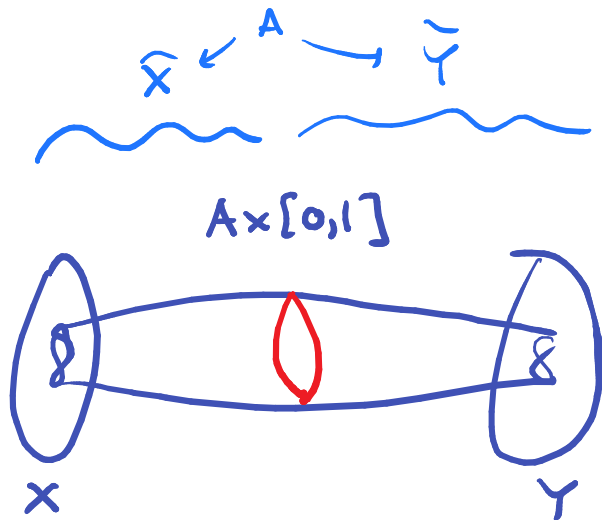
To "compute" homotopy pushout:



Get nice pictures:



\rightsquigarrow



"double mapping cylinder"

homotopy pushout

Homotopy limits/colimits

are "derived functors":

best homotopy invariant approximation to \lim/colim
in Top/sSet
(or in Q.M.C.).

\Rightarrow In ∞ -categories (\Leftrightarrow) "limits/colimits")

Spaces are special: (homotopy theory of Top, or sSet \equiv S)

Homotopy limits/colimits here have additional properties, not satisfied in general. What are these?

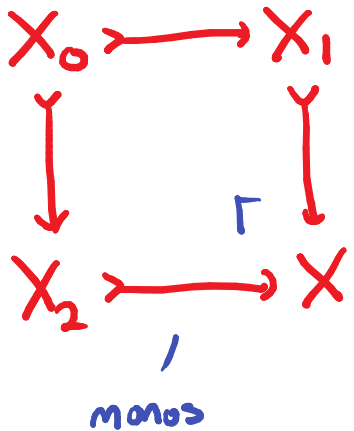
\Rightarrow " ∞ -topos" arises from one answer to this

Set \rightsquigarrow topos

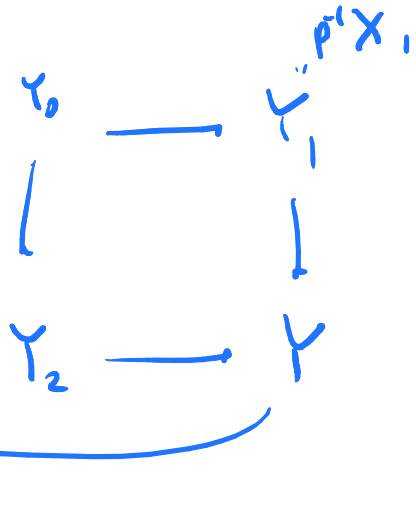
S \rightsquigarrow ∞ -topos

Universality of (homotopy) pushouts: in sSet

homotopy pushout is
always equiv to



in sSet
cof=mono



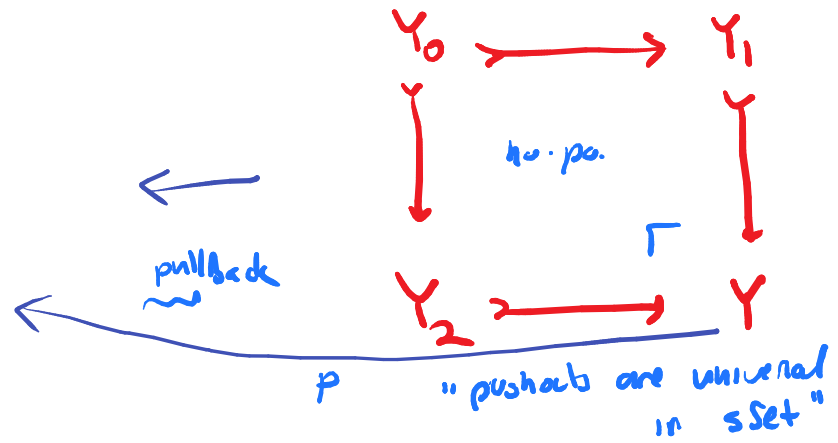
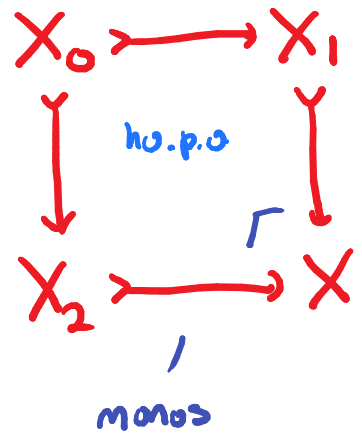
monos

p

Universality of (homotopy) pushouts: in sSet

homotopy pushout is always equiv to

$$Y_1 := X_1 \times_X Y$$



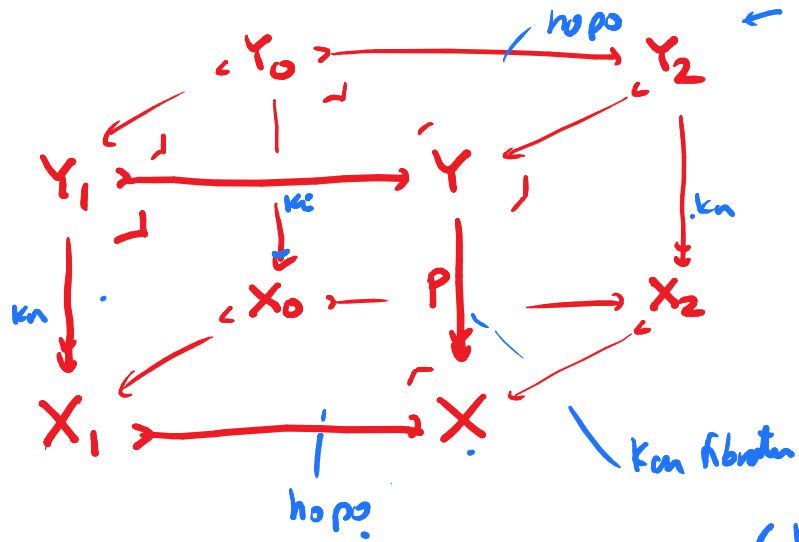
In any 1-topos \mathcal{E} (e.g. \mathbf{sSet}):

colimits are universal: $\forall f: Y \rightarrow X$ in \mathcal{E} :

$f^*: \mathcal{E}_{/X} \rightarrow \mathcal{E}_{/Y}$ preserves colimits

(have right adjoint $\Pi_f: \mathcal{E}_{/Y} \rightarrow \mathcal{E}_{/X}$)

If $p: Y \rightarrow X$ is also a Kan fibration,



all 4 sides are ho pbs.

"ho pushouts are universal w/ homotopy theory of sets"

(also ho. colimit)

Descent for homotopy pushouts:

"opposite composition of operations"

$$\begin{array}{ccccc}
 Y_1 & \longleftarrow & Y_0 & \longrightarrow & Y_2 \\
 \downarrow \text{hopb} & & \downarrow \text{hopb} & & \downarrow \\
 X_1 & \longleftarrow & X_0 & \longrightarrow & X_2
 \end{array}$$

ho colim
 \Longrightarrow

\Longrightarrow
 ho colim

$$\begin{array}{c}
 Y \\
 \downarrow \\
 X
 \end{array}$$



Descent:

$$\begin{array}{ccc}
 Y_i & \longrightarrow & Y \\
 \downarrow & & \downarrow \\
 X_i & \longrightarrow & X
 \end{array}$$

is a homotopy pull back
 $i=0,1,2.$

Example:

$$f: X \rightarrow X$$

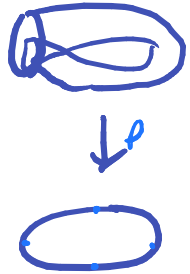
(in Top)
(or sSet)

"mapping cylinder"

$$\begin{array}{ccccc}
 X & \xleftarrow{(id, id)} & X \amalg X & \xrightarrow{(id, f)} & X \\
 \downarrow ho pb & & \downarrow ho pb & & \downarrow ho pb \\
 * & \xleftarrow{\{1, 2\}} & * & \xrightarrow{\{1, 2\}} & *
 \end{array}$$

ho po \Rightarrow

$$C_f \downarrow S^1$$

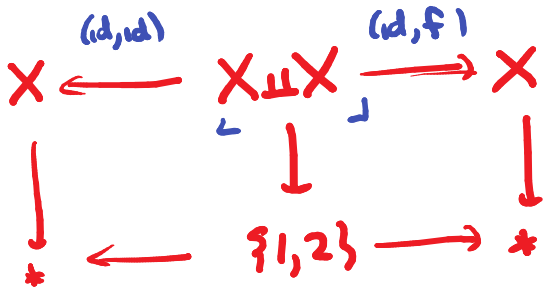


$$C_f = X \times [0, 1] / (x, 1) \sim (fx, 0)$$

f homeomorphism : p fiber bundle, fibers $\cong X$

f homotopy equiv: not fiber bundle, but ho fibers $\cong_{we} X$

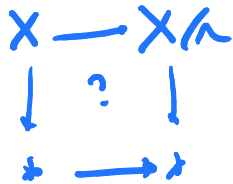
Not true in Set (or any Gr topos): $f: X \xrightarrow[\cong]{\cong} X$



$$\begin{array}{c}
 \text{colim} \\
 \Longrightarrow \\
 X/x \cong f(x)
 \end{array}$$

$$\begin{array}{c}
 \text{colim} \\
 \Longrightarrow \\
 *
 \end{array}$$

\Rightarrow $\tau=1, 2$



p.b only if $f = 1_X$

[can work if
e.g. $\leftarrow x \rightarrow$]

Earliest statement I can find is in

G. Segal. "Categories and cohomology theories" (1974, preprint 1969)
"well-known"

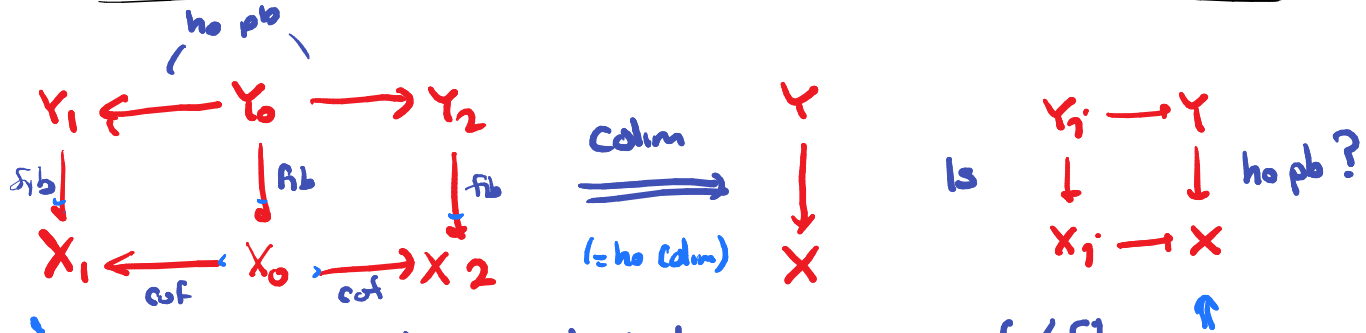
proof sketch in

Puppe, "Remark on 'homotopy fibrations'" (1974)

Independent proof in

Mather, "Pullbacks in homotopy theory" (1976)

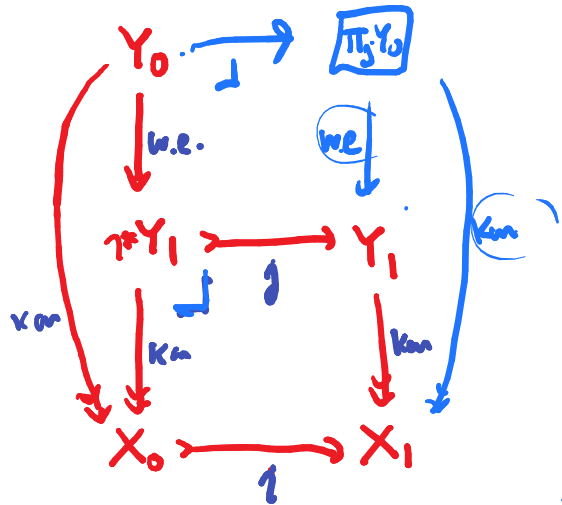
Sketch of proof in sSet model of homotopy theory



Reduce to case where indicated maps are cof/fib

Special case which works: squares ho.pb and pb

Factor homotopy pullback through one which is actual pullback:
(along mono)



Recall : $j : A \rightarrow B$ in $sSet$

$$\Rightarrow sSet_B \begin{array}{c} \xleftarrow{\epsilon_j} \\ \xrightarrow{j^*} \\ \xleftarrow{\pi_j} \end{array} sSet_A$$

$$j \text{ mono} \Rightarrow j^* \circ \pi_j \geq id$$

Miracle : works!

Kapulkin-Lumsdaine (Voevodsky) "The Simplicial model of Univalent Foundations"

Universality of colimits in ∞ -category \mathcal{E}
(cocomplete, finite complete):

For all morphisms $f: X \rightarrow Y$ induced pullback

$f^*: \mathcal{E}/Y \rightarrow \mathcal{E}/X$ preserves colimits

Natural transformation $\lambda: F \Rightarrow F'$ of $F, F': I \rightarrow \mathcal{E}$

is Cartesian iff $\forall \alpha: i \rightarrow j$ in I ,

$$\begin{array}{ccc} F(i) & \xrightarrow{\lambda} & F'(i) \\ \alpha \downarrow & \lrcorner & \downarrow \alpha \\ F(j) & \xrightarrow{\lambda} & F'(j) \end{array}$$

pullback in \mathcal{E} .

$$\mathcal{E}^{\rightarrow} := \text{Fun}(\Delta^1, \mathcal{E})$$

(→)

"arrow category"

$$\text{Cart}(\mathcal{E}^{\rightarrow}) \hookrightarrow \mathcal{E}^{\rightarrow}$$

subcategory (wide, not full)
of Cartesian transformations
 $\downarrow \cong \downarrow$

\mathcal{E} has descent if

- $\text{Cart}(\mathcal{E}^{\rightarrow})$ has colimits
- $\text{Cart}(\mathcal{E}^{\rightarrow}) \rightarrow \mathcal{E}^{\rightarrow}$ preserves colimits

Puppe (1974).

Theorem: Every ∞ -topos has universal colimits + descent

Proof: (1) S \checkmark (uses sSet model) .

(2) $\rightarrow \text{Psh}(C) \hat{=} \text{Fun}(C^{\text{op}}, S)$ lim/colim
computed pointwise

(3) $\text{Psh}(C) \begin{array}{c} \xrightarrow{\ell} \\ \xleftarrow{\gamma} \end{array} E$ ℓ, γ preserve
finite lim / all colims

Theorem: An ∞ -category \mathcal{E} is an ∞ -topos iff

- \mathcal{E} is presentable .
- colimits are universal .
- descent \leftarrow

Sketch \Leftarrow

(1) Find essentially small $\mathcal{C} \in \mathcal{E}$,

• closed under finite limits, st.

$\mathcal{E} \xrightarrow{?} \text{Fun}(\mathcal{C}^{\text{op}}, \mathcal{S})$ is • fully faithful
restricted yoneda • has $\ell \rightarrow i$, ^{and} (i accessible.)
=

application of "accessible ∞ -categories" (Lurie)

e.g. $\mathcal{C} = \mathcal{E}^k \subseteq \mathcal{E}$

↑ "k-compact" objects
("k-presentable")

(2)

Prop: \mathcal{E} - cocomplete, finite complete ∞ -category
universal colimits and descent.

\mathcal{C} - small, finite complete.

$$\mathcal{C} \xrightarrow[\text{Yoneda}]{\rho} \text{Psh}(\mathcal{C}) \xrightarrow{\ell} \mathcal{E}, \quad \ell \text{ colim preserving}$$

Then

$$\ell \text{ lex} \iff \ell \circ \rho \text{ lex}$$

also true, $\mathcal{E} = \text{1-topos}$

$$\text{Fun}(\mathcal{C}^{\text{op}}, \text{Set}) \xrightarrow{\ell} \mathcal{E}$$

Follows from:

Prop': \mathcal{E} - cocomplete, finite complete ω -category
universal colimits and descent.

\mathcal{C} - small

(not finite complete)

$$\mathcal{C} \xrightarrow{P} \text{Psh}(\mathcal{C}) \xrightarrow{L} \mathcal{E}, \quad L \text{ colim preserving}$$

Then

(1) $L(*) \cong *$

$$L \text{ lex} \iff$$

(2)

L preserves pullbacks of form

$$\begin{array}{ccc} \mathcal{B} & \rightarrow & P(\mathcal{C}_2) \\ \downarrow & \searrow & \downarrow \\ P(\mathcal{C}_1) & \rightarrow & P(\mathcal{C}_0) \end{array} \quad \checkmark$$

Lemma: \mathcal{E} - cocomplete, finite complete ω -category
universal colimits
 \mathcal{C} - small

$$\mathcal{C} \xrightarrow{p} \text{Psh}(\mathcal{C}) \xrightarrow{l} \mathcal{E}, \quad l \text{ colim preserving}$$

Then

$$l(X \times Y) \cong lX \times lY \quad \text{iff} \quad l(p(\mathcal{C}) \times p(\mathcal{C}')) = l(p(\mathcal{C})) \times l(p(\mathcal{C}'))$$

(also works if \mathcal{E} has: $\text{colim}_i X_i \times \text{colim}_j Y_j \cong \text{colim}_{i,j} X_i \times Y_j$)

Special Case :

$$l: \text{PSh}(C) \rightarrow \mathcal{E} \quad \text{colim preserving}$$

$$\begin{array}{ccc}
 P & \xrightarrow{\sim} & X' \\
 \downarrow & & \downarrow \\
 X & \xrightarrow{\quad} & B
 \end{array}
 \in \text{Fun}(I, \text{PSh}(C))$$

$\xRightarrow{l} \text{p.b. ?}$

$B := p(C)$

Same as :

$$\text{PSh}(C/c) \simeq \text{prod of rep}$$

product in \Downarrow

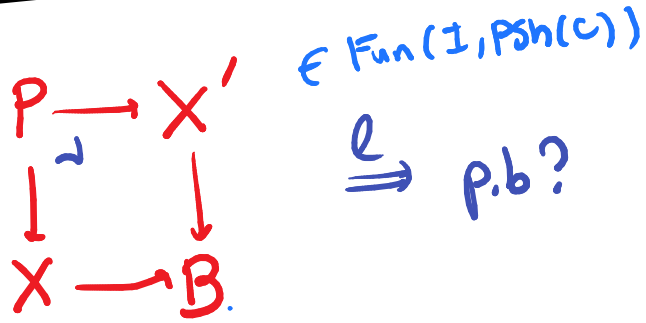
$$\begin{array}{ccc}
 \text{PSh}(C) & / & B \\
 \downarrow p & \xrightarrow{\quad} & \downarrow p(C) \\
 p(C) & \rightarrow & p(C)
 \end{array}$$

$$\xrightarrow{l_B} \mathcal{E}/l(B)$$

\leftarrow Lemma

have univ colim, descent
pres products

General case: $\ell: \text{Psh}(C) \rightarrow \mathcal{E}$ colim preserving

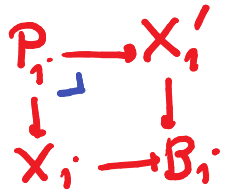


$\in \text{Fun}(I, \text{Psh}(C))$
 $\xRightarrow{\ell} \text{p.b.}?$

Write $B \cong \text{colim}_{i \in I} B_i$,

B_i representable.
 $\in \text{Psh}(C)$

\Rightarrow Pull back to
 (along $B_i \rightarrow B$)



$i \in I$

$$\begin{array}{ccc}
 P \rightarrow X' & \xleftarrow{\langle \text{Func}, \text{Psh}(C) \rangle} & P_i \rightarrow X_i' \\
 \downarrow \dashv \quad \downarrow & \rightsquigarrow & \downarrow \dashv \quad \downarrow \\
 X \rightarrow B & & X_i \rightarrow B_i
 \end{array}
 \quad \begin{array}{l}
 \text{Psh}(C) \\
 B = \text{colim } B_i \implies \\
 \leftarrow \text{rep}
 \end{array}$$

$$\begin{array}{ll}
 P = \text{colim } P_i & \text{colims are} \\
 X = \text{colim } X_i & \text{univ in} \\
 X' = \text{colim } X_i' & \text{Psh}(C)
 \end{array}$$

and $\forall i \rightarrow j$ in I

$$\begin{array}{ccc}
 P_i \rightarrow X_i \rightarrow B_i & & P_i \rightarrow X_i' \rightarrow B_i \\
 \downarrow \dashv \quad \downarrow \dashv \quad \downarrow & & \downarrow \dashv \quad \downarrow \dashv \quad \downarrow \\
 P_j \rightarrow X_j \rightarrow B_j & \leftarrow \text{rep} & P_j \rightarrow X_j' \rightarrow B_j \leftarrow \text{rep}
 \end{array}$$

\nearrow ps in Psh(C)

\mathcal{L} preserves these
pullbacks over a
representable
(special case)

$$\begin{array}{ccc}
 \mathbb{L}P_i \rightarrow \mathbb{L}X_i \rightarrow \mathbb{L}B_i \\
 \downarrow \quad \downarrow \quad \downarrow \\
 \mathbb{L}P_j \rightarrow \mathbb{L}X_j \rightarrow \mathbb{L}B_j
 \end{array}$$

\uparrow \uparrow
 Cart rest frnd

and

$$\begin{aligned}
 \mathbb{L}P &= \text{colim } \mathbb{L}P_i \\
 \mathbb{L}X &= \text{colim } \mathbb{L}X_i \\
 \mathbb{L}B &= \text{colim } \mathbb{L}B_i
 \end{aligned}$$

descent

$$\begin{array}{ccc}
 \mathbb{L}P_i \rightarrow \mathbb{L}X_i \rightarrow \mathbb{L}B_i \\
 \downarrow \quad \downarrow \quad \downarrow \\
 \mathbb{L}P \rightarrow \mathbb{L}X \rightarrow \mathbb{L}B
 \end{array}$$

\mathbb{L} colim

$$I \rightarrow \text{Cart}(\mathcal{E} \rightarrow) \rightarrow \mathcal{E} \rightarrow$$

pullback of

$$\mathbb{L}P \xrightarrow{f} \mathbb{L}X \times_{\mathbb{L}B} \mathbb{L}X'$$

along each $\mathbb{L}B_i \rightarrow \mathbb{L}B$ is iso

\Rightarrow univ colm in $\mathcal{E} \Rightarrow f$ iso

$$\mathbb{L}P_i \rightarrow \mathbb{L}X_i \xrightarrow{f_i} \mathbb{L}X'_i$$

Next time:

• "Object classifier" ←

• Truncation + Connectivity ←